

Noisy nets and transient limits

Fred Richman

Florida Atlantic University

29 August 2019

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Paul Halmos \longrightarrow Errett Bishop \longrightarrow Y.K. Chan

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Schizophrenia in contemporary mathematics, August 1973

The motivating example

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Our interest is in long-term collective opinion; in particular, the possible persistence of a minority opinion.

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$$\begin{array}{cccccc}
 1 & 1 & 0 & 0 & 0 & 0 \\
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Alternatively, we could posit that the configuration $\blacksquare\blacksquare\square\square\blacksquare\square$ is a circle—the first individual being adjacent to the last. In that case the state would switch to $\blacksquare\blacksquare\square\square\square\blacksquare$, which is where it would stay.

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For $m = 4$ these are exactly the six strings

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There is also a pair of strings of period two (only for m even):

$$\blacksquare\square\blacksquare\square \longleftrightarrow \square\blacksquare\square\blacksquare$$

The 12 fixed points for the 5-circle are



and their complements.

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and 2206 fixed points for the 16-circle.

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The fixed states are where each square is next to two squares of the same color.

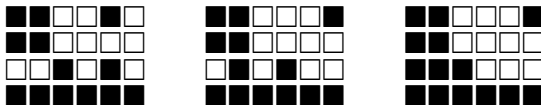
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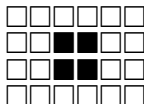
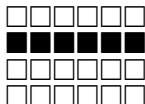
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Here's a 4-by-6 example and how it steps to a fixed point:

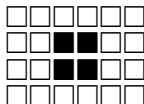
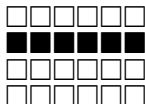


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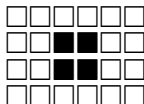
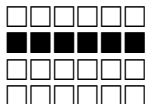


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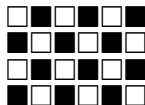
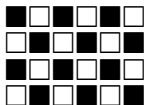


and a pair of points of period two:

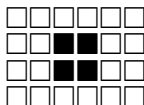
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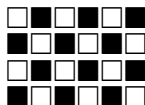
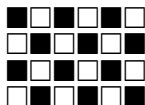
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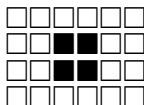


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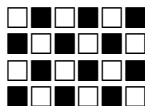
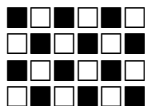


If we change these states of period two by filling in all four squares in the upper right corner, we get another pair of states of period two:

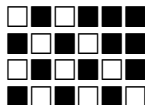
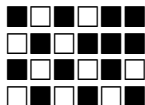
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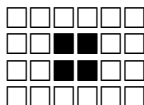
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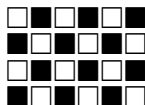
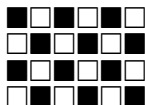
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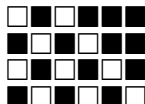
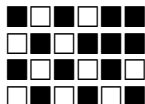
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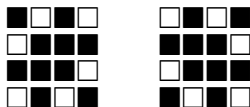
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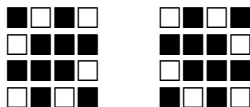
That's essentially the same as filling in the four central squares.

We can do the same in the 4-by-4 torus:

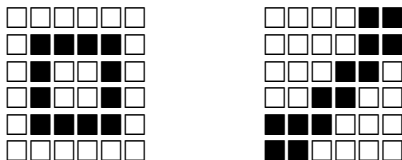
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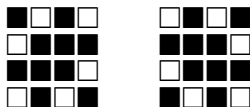
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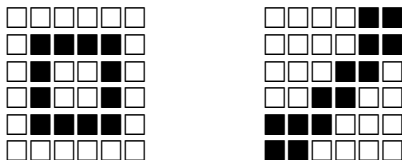
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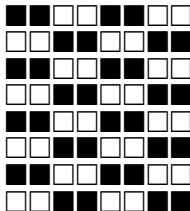
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and a point of period two of the 8-by-8 torus:



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for each $\theta > 0$.

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for each $\theta > 0$. Like the standard normal density function

$$f(x) = \frac{1}{\sqrt{2\pi}} e^{-x^2/2}$$

Noise

Introduce noise by changing the stepping equation from $v = \text{sgn } Cu$ to $v = \text{sgn } (Cu + X)$, where X is a vector of independent random variables with the same (parameterized) density function.

Start with a continuous density function f that is everywhere positive, symmetric around 0, and satisfies

$$\lim_{x \rightarrow \infty} \frac{f(\theta x)}{x} = 0$$

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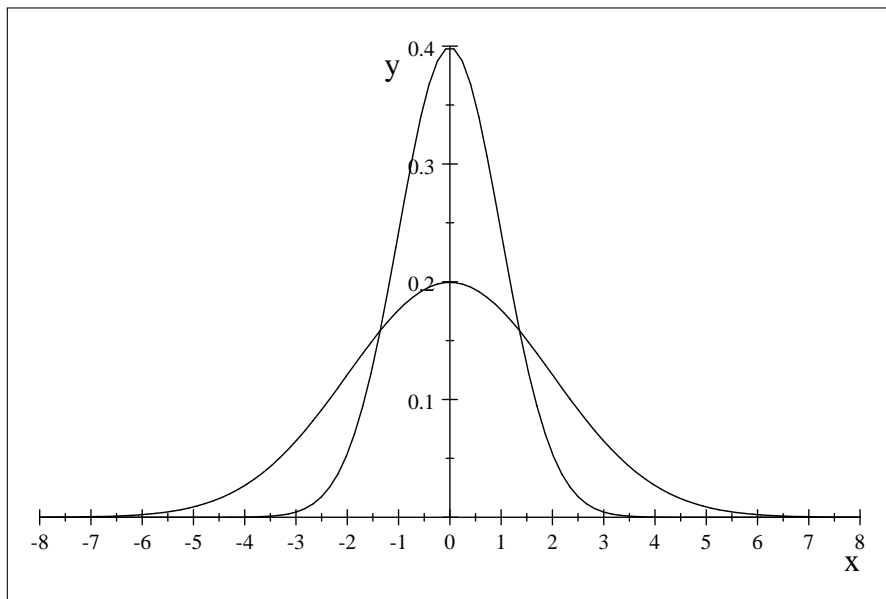
$$f(x) = \frac{1}{\sqrt{2\pi}} e^{-x^2/2}$$

Then pass to the parameterized density function

$$\frac{1}{\sigma} f\left(\frac{x}{\sigma}\right)$$

Here are graphs of the Gaussian density function for $\sigma = 1$ and $\sigma = 2$.

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For the Gaussian distribution, $[4](1) = 0.00003167$ and $[4](2) = 0.02275$.

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	000	001	010	011	100	101	110	111
000	$\bar{2}\bar{3}\bar{2}$	$\bar{2}\bar{3}2$	$\bar{2}3\bar{2}$	$\bar{2}32$	$2\bar{3}\bar{2}$	$2\bar{3}2$	$23\bar{2}$	232
001	$\bar{2}\bar{1}0$	$\bar{2}\bar{1}0$	$\bar{2}10$	$\bar{2}10$	$2\bar{1}0$	$2\bar{1}0$	210	210
010	$0\bar{2}0$	$0\bar{2}0$	020	020	$0\bar{2}0$	$0\bar{2}0$	020	020
011	$01\bar{2}$	$01\bar{2}$	$0\bar{1}2$	$0\bar{1}2$	012	$01\bar{2}$	$0\bar{1}2$	$0\bar{1}2$
100	$0\bar{1}\bar{2}$	$0\bar{1}\bar{2}$	$01\bar{2}$	012	$0\bar{1}\bar{2}$	$0\bar{1}\bar{2}$	$01\bar{2}$	012
101	020	020	$0\bar{2}0$	$0\bar{2}0$	020	020	$0\bar{2}0$	$0\bar{2}0$
110	210	210	$2\bar{1}0$	$2\bar{1}0$	$\bar{2}10$	$\bar{2}10$	$\bar{2}\bar{1}0$	$\bar{2}\bar{1}0$
111	232	$23\bar{2}$	$2\bar{3}2$	$2\bar{3}\bar{2}$	$\bar{2}32$	$\bar{2}3\bar{2}$	$\bar{2}\bar{3}2$	$\bar{2}\bar{3}\bar{2}$

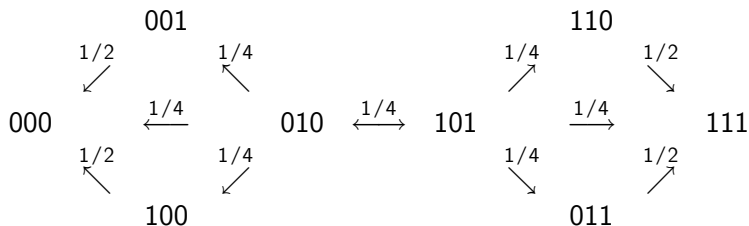
Here is the matrix $A(0)$ for the three-node line.

	000	001	010	011	100	101	110	111
000	1	0	0	0	0	0	0	0
001	1/2	1/2	0	0	0	0	0	0
010	1/4	1/4	0	0	1/4	1/4	0	0
011	0	0	0	1/2	0	0	0	1/2
100	1/2	0	0	0	1/2	0	0	0
101	0	0	1/4	1/4	0	0	1/4	1/4
110	0	0	0	0	0	0	1/2	1/2
111	0	0	0	0	0	0	0	1

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011	0	0	0	1/2	0	0	0	1/2
100	1/2	0	0	0	1/2	0	0	0
101	0	0	1/4	1/4	0	0	1/4	1/4
110	0	0	0	0	0	0	1/2	1/2
111	0	0	0	0	0	0	0	1

It represents the Markov chain



In the limit, this Markov chain is $000 \xleftarrow{3/4} 010 \xleftrightarrow{1/4} 101 \xrightarrow{3/4} 111$ so 010 becomes 000 with probability

$$3/4 (1 + 1/16 + 1/16^2 + \dots) = \frac{3 \cdot 16}{4(16 - 1)} = \frac{4}{5}$$

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	000	001	010	011	100	101	110	111
000	1	0	0	0	0	0	0	0
001	1	0	0	0	0	0	0	0
010	4/5	0	0	0	0	0	0	1/5
011	0	0	0	0	0	0	0	1
100	1	0	0	0	0	0	0	0
101	1/5	0	0	0	0	0	0	4/5
110	0	0	0	0	0	0	0	1
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100	1	0	0	0	0	0	0	0
101	1/5	0	0	0	0	0	0	4/5
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111	0	0	0	0	0	0	0	1

This is actually a transient limit for powers of $A(\sigma)$.

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A **window** is a pair of real valued functions of a positive real variable σ such that

$$1 \leq u(\sigma) \leq v(\sigma)$$

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If $A(\sigma)$ is an indexed (Markov) matrix, then L is a **transient limit** of the sequence $A^n(\sigma)$ if there is a window $u \leq v$ such that for $u \leq n \leq v$, the matrices $A^n(\sigma)$ and L are close for all sufficiently small values of σ .

Example 1. Let A be the 2-by-2 Markov matrix $\begin{pmatrix} 1 - \sigma & \sigma \\ 0 & 1 \end{pmatrix}$.

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Note that $\lim_{\sigma \rightarrow 0} (1-\sigma)^{1/\sqrt{\sigma}} = 1$ and $\lim_{\sigma \rightarrow 0} (1-\sigma)^{1/\sigma^2} = 0$.

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The s_i , for $i = 1, 2, \dots, m - 1$, are transition probabilities (indexed by σ) such that $\lim_{\sigma \rightarrow 0} s_i / s_{i-1} = 0$ where $s_0 = 1$.

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The transition matrix A is

$$\begin{pmatrix} 1 - s_1 & s_1 & 0 & \dots & 0 & 0 \\ 0 & 1 - s_2 & s_2 & \dots & 0 & 0 \\ 0 & 0 & 1 - s_3 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & 1 - s_{m-1} & s_{m-1} \\ 0 & 0 & 0 & \dots & 0 & 1 \end{pmatrix}$$

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In the window defined by $u = 1$ and $v = \sqrt{1/[1](\sigma)}$, the transient limit is the matrix

	000	001	010	011	100	101	110	111
000	1	0	0	0	0	0	0	0
001	1	0	0	0	0	0	0	0
010	4/5	0	0	0	0	0	0	1/5
011	0	0	0	0	0	0	0	1
100	1	0	0	0	0	0	0	0
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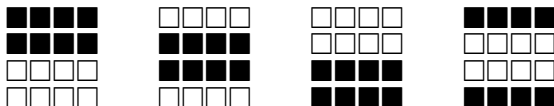
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Example 4. The 4-by-4 torus. Let $\tau = \frac{1}{[2](\sigma)}$ and consider the window $u = \tau$ and $v = \tau^2$.

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Yay!

Jessica Su